Optimization Techniques in Finance
2. Constraint optimization and Lagrange multipliers

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Fall 2019
Outline

1. Constraint optimization problems
2. Numerical methods
In most financial applications, the variables in an optimization problem are restricted to vary in a subset $\Omega$ of $\mathbb{R}^n$ rather than in the entire space $\mathbb{R}^n$.

Examples:

(i) The correlation parameter in a stochastic volatility model has to lie within the range $[-1, 1]$.

(ii) The volatility parameters in a term structure model should be all positive, while the unconstrained optimizer stubbornly produces negative values.

(iii) Portfolio manager’s mandate is to be long only, meaning that the weights of all assets in the portfolio are positive.

(iv) The share of each position in the portfolio cannot exceed 5% of its total value.

(v) No industry group should take more than 20% share of the total portfolio value.

Oftentimes we are “lucky” and get the right results using an unconstrained method (like in data fitting). Otherwise, we have to resort to the specialized methods developed to handle constrained problems.
We shall assume that the subset $\Omega$ is specified in terms of a number of functions $c_i(x)$, $i = 1, \ldots, m$, which define equality constraints or inequality constraints.

Specifically, given an objective function $f(x)$, we formulate the following problem:

$$\min_{x \in \mathbb{R}^n} f(x), \quad \text{subject to } \begin{cases} c_i(x) = 0, & \text{if } i \in \mathcal{E}, \\ c_i(x) \leq 0, & \text{if } i \in \mathcal{I}. \end{cases}$$

(1)

Here, $\mathcal{E}$ and $\mathcal{I}$ are disjoint subsets of the set of indices $1, \ldots, m$, such that $\mathcal{E} \cup \mathcal{I} = \{1, \ldots, m\}$.

A point $x \in \mathbb{R}^n$ is called feasible, if it satisfies all the constraints. We can thus characterize the subset $\Omega$ as the set of all feasible points of the problem:

$$\Omega = \{x \in \mathbb{R}^n : c_i(x) = 0, \text{if } i \in \mathcal{E}; c_i(x) \leq 0, \text{if } i \in \mathcal{I}\}.$$  

(2)

A point $x^*$ is a local solution to (1), if it is feasible, and $f(x) \geq f(x^*)$, for all $x$ in a neighborhood of $x^*$ (which is contained in $\Omega$).

It is a strict local solution if $f(x) > f(x^*)$, for all $x \neq x^*$ in a neighborhood of $x^*$. 
Solving constraint optimization problems is challenging and, before developing general methodology, we discuss a few examples.

**Example 1.** Consider the problem:

$$\min x_1 + x_2, \quad \text{subject to } \begin{cases} x_1 x_2 = a^2, \text{ where } a > 0, \\ x_1, x_2 \geq 0. \end{cases} \quad (3)$$

In this case, $c_1(x) = x_1 x_2 - a^2$, $\mathcal{E} = \{1\}$, $c_2(x) = -x_1$, $c_3(x) = -x_2$, $\mathcal{I} = \{2, 3\}$, and the feasible set is the hyperbola $x_1 x_2 = a^2$ in the first quadrant.

The special feature of this problem is that the constraints can be solved. Namely, $x_2 = a^2 / x_1$, which reduces the problem to minimizing a function of one variable:

$$g(x_1) = x_1 + \frac{a^2}{x_1}.$$ 

Setting the derivative of $g(x_1)$ to zero we find that $x_1^* = a$ (the solution $-a$ is rejected because it is not in the feasible set). That means that $x_2^* = a$, and inspection shows that $x^* = (a, a)$ is, in fact, a global minimum.

Situations in which the constraints can be solved are rare. In general, solving the constraints is either impossible or it leads to cumbersome calculations.
Example 2. Consider the problem:

\[ \min x_1 + x_2, \quad \text{subject to } x_1^2 + x_2^2 = 2. \]  

(4)

In this case, \( c_1(x) = x_1^2 + x_2^2 - 2 \), \( \mathcal{E} = \{1\} \), \( \mathcal{I} = \emptyset \), and the feasible set consists of the circle of radius \( \sqrt{2} \).

Inspection shows that the solution to this problem is

\[ x^* = -\begin{pmatrix} 1 \\ 1 \end{pmatrix}. \]

We note that

\[ \nabla f(x^*) = -\frac{1}{2} \nabla c_1(x^*), \]

(5)

i.e., at the solution, the gradient of the objective function is proportional to the gradient of the constraint function.
This is not a coincidence! We will see that this is a general fact: at a local minimizer, the gradient of the objective function is a linear combination of the gradients of the constraints.

We rewrite (5) in the form:

\[ \nabla f(x^*) + \frac{1}{2} \nabla c_1(x^*) = 0. \]  

(6)

The proportionality coefficient \( \lambda^* = 1/2 \) (in this example) is called the Lagrange multiplier.

We can interpret this observation geometrically as follows.
If $x$ is a feasible point that *is not* a solution to (4), then there is a small vector $h = \varepsilon d$, $\varepsilon > 0$, such that $x + h$ is feasible and $f(x + h) < f(x)$, i.e.

\[
0 = c_1(x + h) \\
\approx \nabla c_1(x)^T h, \\
0 > f(x + h) - f(x) \\
\approx \nabla f(x)^T h.
\]

Therefore, for $x$ to be a solution, there *cannot* exist a vector (direction) $d$ such that the conditions below hold simultaneously:

\[
\nabla c_1(x)^T d = 0, \\
\nabla f(x)^T d < 0.
\]

Thus $x$ can be a solution if and only if $\nabla f(x)$ and $\nabla c_1(x)$ are parallel. In other words, there has to exist a scalar $-\lambda$ such that $\nabla f(x) = -\lambda \nabla c_1(x)$. 
Let us now consider the general constrained optimization problem with equality constraints only (i.e. $\mathcal{I} = \emptyset$).

Reasoning along the lines of Example 2, we argue that a feasible point $x$ is a solution to (1), provided that there is no vector $d \in \mathbb{R}^n$, $d \neq 0$, with the properties:

$$\nabla c_i(x)^T d = 0, \text{ for } i = 1, \ldots, m,$$

$$\nabla f(x)^T d < 0. \quad (7)$$

We shall call a feasible point $x$ \textit{regular}, if the vectors $\nabla c_1(x), \ldots, \nabla c_m(x)$ are linearly independent.

If $x$ is regular, the $m$ vectors $\nabla c_i(x) \in \mathbb{R}^n$ span an $m$-dimensional subspace $W(x) \subset \mathbb{R}^n$.

The first condition in (7) defines the subspace of \textit{first order feasible variations}:

$$V(x) = \{d : \nabla c_i(x)^T d = 0, \text{ for } i = 1, \ldots, m\}. \quad (8)$$

Notice that $V(x)$ is simply the orthogonal complement of $W(x)$, $V(x) = W(x)\perp$. 

Equality constraints and Lagrange Multiplier Theorem
The (impossibility of the) second condition in (7) implies that the gradient $\nabla f(x)$ of the objective function has to be perpendicular to $V(x)$.

Indeed, if the inner product of $\nabla f(x)$ with a nonzero $d \in V(x)$ is positive then, by replacing $d$ with $-d$, it can be negative, which is assumed impossible. Therefore, $\nabla f(x) \in V(x)^\perp$.

But $V(x)^\perp = W(x)^\perp = W(x)$.

Consequently, $\nabla f(x)$ must be a linear combination of the constraint gradients $\nabla c_i(x)$, which span $W(x)$, i.e.

$$\nabla f(x) + \sum_{i=1}^{m} \lambda_i \nabla c_i(x) = 0,$$

with some scalars $\lambda_i$. 

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The Lagrange Multiplier Theorem formulated below states necessary conditions for local minima of (1). It puts the informal reasoning above on a rigorous basis.

**Lagrange Multiplier Theorem.** Let $x^*$ be a regular local minimizer of $f(x)$ subject to $c_i(x) = 0$, for $i = 1, \ldots, m$. Then:

(i) There exists a unique vector $\lambda^* = (\lambda_1^*, \ldots, \lambda_m^*)$ of Lagrange multipliers, such that

$$
\nabla f(x^*) + \sum_{i=1}^{m} \lambda_i^* \nabla c_i(x^*) = 0.
$$

(ii) If, in addition, $f(x)$ and $c_i(x)$ are twice continuously differentiable, then

$$
d^\top \left( \nabla^2 f(x^*) + \sum_{i=1}^{m} \lambda_i^* \nabla^2 c_i(x^*) \right) d \geq 0,
$$

for all $d \in V(x^*)$. 

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It is convenient (and customary) to write the necessary conditions in terms of the Lagrangian function. This is a function of \( n + m \) variables defined as follows:

\[
L(x, \lambda) = f(x) + \sum_{i=1}^{m} \lambda_i c_i(x).
\]  

(11)

The necessary condition of the Lagrange Multiplier Theorem can be parsimoniously formulated as follows:

\[
\nabla_x L(x^*, \lambda^*) = 0, \\
\nabla_\lambda L(x^*, \lambda^*) = 0, \\
d^T \nabla^2_{xx} L(x^*, \lambda^*) d \geq 0, \text{ for all } d \in V(x^*).
\]  

(12)

We emphasize that these conditions are necessary but not sufficient: a solution to the system above may not represent a local minimum.
Example 3. Consider the problem:

$$\min \frac{1}{2} (x_1^2 + x_2^2 + x_3^2), \quad \text{subject to } x_1 + x_2 + x_3 = 3.$$

The first order necessary conditions read:

$$x_i + \lambda = 0, \text{ for } i = 1, 2, 3,$$

$$x_1 + x_2 + x_3 = 3.$$

Solving this system yields

$$x^* = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \quad \lambda^* = -1.$$

Note also that the second order (positive definiteness) condition holds, as

$$\nabla_{xx}^2 L(x^*, \lambda^*) = I_3 \text{ (the } 3 \times 3 \text{ identity matrix). In fact, } x^* \text{ is a global minimum.}$$
The theorem below gives a sufficient condition for the existence of a local minimizer.

**Second Order Sufficient Conditions.** Let $f(x)$ and $c_i(x)$ be twice continuously differentiable and let $x^* \in \Omega \subset \mathbb{R}^n$, $\lambda^* \in \mathbb{R}^m$, be such that

\[
\nabla_x L(x^*, \lambda^*) = 0, \\
\nabla_\lambda L(x^*, \lambda^*) = 0,
\]

(13)

and

\[\quad d^T \nabla_{xx}^2 L(x^*, \lambda^*) d > 0, \text{ for all } d \in V(x^*).\]

(14)

Then $x^*$ is a strict local minimizer of $f(x)$ subject to equality constraints $c_i(x), i = 1, \ldots, m.$
Example 4. Consider the quadratic optimization problem:

$$\min f(x) = \frac{1}{2} x^T A x + x^T b, \quad \text{subject to } x^T c = 1, \quad (15)$$

where $b, c \in \mathbb{R}^n$ are constant vectors, and $A \in \text{Mat}_n(\mathbb{R})$ is positive definite.

First and second order conditions yield

$$\lambda^* = -\frac{b^T A^{-1} c + 1}{c^T A^{-1} c},$$

$$x^* = -A^{-1}(b + \lambda^* c),$$

$$\nabla^2_{xx} L(x^*, \lambda^*) = A.$$

Note that, since $A$ is positive definite, condition (14) holds for all $d$ (and in particular, for $d \in V(x^*) = \{d : c^T d = 0\}$).

Consequently, $x^*$ is a strict (global) minimizer.
Lagrange multipliers have often intuitive interpretation, depending on the specific problem at hand.

In general, they can be interpreted as the \textit{rates of change} of the objective function as the constraint functions are varied.

Let \( x^* \) and \( \lambda^* \) be a local minimizer and the corresponding Lagrange multiplier, respectively, of a constrained optimization problem for \( f(x) \).

Consider now the following family of constrained optimization problems, parameterized by a vector \( u = (u_1, \ldots, u_m) \in \mathbb{R}^m \):

\[
\min_{x \in \mathbb{R}^n} f(x), \quad \text{subject to} \quad c_i(x) - u_i = 0. \tag{16}
\]

One can think of the family of constraint functions \( c_i(x, u) = c_i(x) - u_i \) as defining the parametric problem.

Then, at least for small values of \( u \), the solution \( x(u) \) with the corresponding Lagrange multipliers \( \lambda(u) \) forms a continuously differentiable function, such that

\[
x(0) = x^*, \\
\lambda(0) = \lambda^*.
\]
Observe that, since $\nabla_x c_i(x, u) = \nabla_x c_i(x)$, we have the following equation:

$$
\nabla_x f(x(u)) + \sum_{i=1}^{m} \lambda_i(u) \nabla_x c_i(x(u)) = 0,
$$

which is simply the first order condition for (16).

We denote

$$p(u) = f(x(u)),$$

the minimum of the objective function as a function of the parameters $u$.

The function $p(u)$ is called the *primal function*.

Then, for all $u$ where $x(u)$ is defined, we have the following relation:

$$\nabla_u p(u) = -\lambda(u).$$

In words: the Lagrange multipliers are the (negative of the) rates of change of the primal function as the parameters are varied.
The proof is a straightforward calculation. From the chain rule and (17), we get for each $j = 1, \ldots, m$,

$$\nabla_{u_j} p(u) = \nabla_{u_j} f(x(u))$$

$$= \sum_{k=1}^{n} \nabla_{x_k} f(x(u)) \nabla_{u_j} x_k(u)$$

$$= - \sum_{i=1}^{m} \sum_{k=1}^{n} \lambda_i(u) \nabla_{x_k} c_i(x(u)) \nabla_{u_j} x_k(u).$$

However, using the chain rule again,

$$\sum_{k=1}^{n} \nabla_{x_k} c_i(x(u)) \nabla_{u_j} x_k(u) = \nabla_{u_j} c_i(x(u))$$

$$= \nabla_{u_j} u_i$$

$$= \delta_{ji},$$

where we have also used the fact that $c_i(u) = u_i$. 
As a consequence,

\[ \nabla_u \rho(u) = - \sum_{i=1}^{m} \lambda_i(u) \delta_{ji} \]

\[ = - \lambda_j(u), \]

which proves equation (19).
We will now move on to the case of inequality constraints. We begin with a motivational example.

**Example 5.** Consider the slightly modified Example 2:

\[
\min x_1 + x_2, \quad \text{subject to } x_1^2 + x_2^2 \leq 2.
\]

(20)

In this case, \( c_1(x) = x_1^2 + x_2^2 - 2 \), \( \mathcal{E} = \emptyset \), \( \mathcal{I} = \{1\} \), and the feasible set consists of the unit circle and its interior.

Inspection shows that the solution to this problem continues to be

\[
x^* = -\begin{pmatrix} 1 \\ 1 \end{pmatrix}.
\]

Notice that condition (5) continues to hold. We will argue that, in case of an inequality constraint, the *sign* of the Lagrange multiplier is not a coincidence.
Assume that a feasible point $x \in \mathbb{R}^2$ is not a local minimizer.

Arguing along the lines of Example 2, this is possible, if there cannot exist a (small) vector $h = \varepsilon d$, $\varepsilon > 0$, satisfying both conditions below hold simultaneously:

$$c_1(x) + \varepsilon \nabla c_1(x)^T d \leq 0,$$
$$\nabla f(x)^T d < 0.$$  \hspace{1cm} (21)

We consider two cases.

Case 1. $x$ lies inside the circle, i.e. $c_1(x) < 0$ (the constraint is inactive). In this case, assuming that $\nabla f(x) \neq 0$, the vector $d = -\nabla f(x)$ satisfies both conditions, provided that $\varepsilon > 0$ is sufficiently small.

Note that this vector does not work if $\nabla f(x) = 0$. 

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Case 2. $x$ lies on the boundary of the circle, i.e. $c_1(x) = 0$ (the constraint is active).

In this case, the conditions read

$$\nabla c_1(x)^T d \leq 0,$$
$$\nabla f(x)^T d < 0.$$ (22)

The first of these conditions defines a closed half-plane, while the second one defines an open half-plane.

The intersection of these two half-planes should be empty!

A reflection shows that this is possible only if there is a positive constant $\lambda$ such that

$$\nabla f(x) = -\lambda \nabla c_1(x).$$
Inequality constraints

- We can formulate the results of the analysis of Cases 1 and 2 in the following elegant way using the Lagrange function.

  - If $x^*$ is a local minimizer of $f(x)$ (no feasible descent direction $d$ is possible), then
    
    $$\nabla_x L(x^*, \lambda^*) = 0, \text{ for some } \lambda^* \geq 0,$$

    with the additional requirement that

    $$\lambda^* c_1(x^*) = 0.$$

  - The latter condition is called the **complementary slackness condition**. It means that $\lambda^*$ can be strictly positive, only if the constraint $c_1(x)$ is active.

  - In Case 1 $c_1(x^*) < 0$, and so $\lambda^* = 0$, while in Case 2 $\lambda^* = 1/2$ is positive.
Consider now the general problem (1), in which both equality and inequality constraints are present.

For any feasible point \( x \in \Omega \), we define the set of active inequality constraints by

\[
A(x) = \{ i \in I : c_i(x) = 0 \}.
\]  

(23)

A constraint \( i \in I \) for which \( c_i(x) < 0 \) is called *inactive* at \( x \).

If \( x^* \) is a local minimum solution to (1), then it is also a local minimum to the same problem with all constraints inactive at \( x^* \) ignored. Thus the inactive constraints do not matter.

The active constraints can be treated as equality constraints. A local minimum solution to (1) solves thus the equality constraints problem:

\[
\min_{x \in \mathbb{R}^n} f(x), \quad \text{subject to } c_i(x) = 0, \text{ if } i \in \mathcal{E} \cup A(x).
\]  

(24)
Consequently, assuming that $x^*$ is regular, there exist Lagrange multipliers $\lambda_i^*, i \in \mathcal{E} \cup \mathcal{A}(x^*)$, so that

$$\nabla f(x^*) + \sum_{i \in \mathcal{E} \cup \mathcal{A}(x^*)} \lambda_i^* \nabla c_i(x^*) = 0.$$  

For convenience, we also introduce the Lagrange multiplier $\lambda_i^* = 0$ for each inactive constraint $i$. We can thus write the above condition compactly as

$$\nabla f(x^*) + \sum_{i=1}^{m} \lambda_i^* \nabla c_i(x^*) = 0,$$

$$\lambda_i^* c_i(x^*) = 0.$$  

An important fact about the Lagrange multipliers corresponding to the inequality constraints is that they are nonnegative.
This can be seen as follows.

If the $i$-th constraint is replaced by a weaker constraint $c_i(x) \leq u_i$, where $u_i > 0$, then the optimal objective function will not increase. The reason for this is that the feasible will become larger.

From the sensitivity equation (19):

$$\lambda_i^* = -\lim_{u_i \downarrow 0} \frac{\partial}{\partial u_i} p(u) \geq 0.$$  

The arguments above are not exactly a proof, but are convincing enough to understand the following necessary condition.
As in the case of equality constraints only, it is convenient to work with the Lagrangian function:

\[ L(x, \lambda) = f(x) + \sum_{i=1}^{m} \lambda_i c_i(x). \]  \hspace{1cm} (26)

**Karush-Kuhn-Tucker Necessary Conditions.** Let \( x^* \) be a local minimum solution to (1), and assume that \( x^* \) is regular. Then there exists a unique vector of Lagrange multipliers \( \lambda_i^*, \ i = 1, \ldots m \), such that

\[ \nabla_x L(x^*, \lambda^*) = 0, \]
\[ \lambda_i^* c_i(x^*) = 0, \]
\[ \lambda_i^* \geq 0, \text{ for } i \in \mathcal{I}. \]  \hspace{1cm} (27)

If, additionally, \( f(x) \) and \( c_i(x) \) are twice continuously differentiable, then

\[ d^T \nabla_x^2 L(x^*, \lambda^*) d \geq 0, \]  \hspace{1cm} (28)

for all \( d \), such that \( \nabla c_i(x^*)^T d = 0, \ i \in \mathcal{E} \cup \mathcal{A}(x^*). \)
Inequality constraints and the KKT conditions

Note that the complementary slackness condition

$$\lambda_i^* c_i(x^*) = 0$$  \hspace{1cm} (29)

is just a compact way of stating the fact $\lambda_i^* = 0$, if $c_i(x^*)$ is not active, and it can be nonzero (and, in fact, nonnegative), if the constraint is active.
Example 6. Consider the problem:

$$\min \frac{1}{2} (x_1^2 + x_2^2 + x_3^2), \quad \text{subject to } x_1 + x_2 + x_3 \leq -3.$$

The first order necessary conditions read:

$$x_i + \lambda = 0, \quad \text{for } i = 1, 2, 3,$$

or

$$x_i = -\lambda, \quad \text{for } i = 1, 2, 3.$$
We now have two possibilities:

(i) The constraint is inactive, \( x_1^* + x_2^* + x_3^* < -3 \), in which case \( \lambda^* = 0 \). This implies \( x_i^* = 0 \), for \( i = 1, 2, 3 \), which contradicts the constraint. Thus \((0, 0, 0)\) is not a minimizer.

(ii) The constraint is active, \( x_1^* + x_2^* + x_3^* = -3 \), in which case \( \lambda^* = 1 \), and

\[
x^* = -\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}.
\]

Note also that the second order (positive definiteness) condition holds, as

\[
\nabla_x^2 L(x^*, \lambda^*) = I_3.
\]

A reflection shows that \( x^* \) is, in fact, a global minimizer.
Example 7. Consider the problem:

$$
\min 2(x_1 - 1)^2 + (x_2 - 2)^2,
$$
subject to

$$
\begin{align*}
& x_1 + 2x_2 \leq 1, \\
& x_1 \geq x_2.
\end{align*}
$$

The Lagrange function is

$$
L(x, \lambda) = 2(x_1 - 1)^2 + (x_2 - 2)^2 + \lambda_1 (x_1 + 2x_2 - 1) + \lambda_2 (-x_1 + x_2),
$$

and the KKT conditions read:

$$
\begin{align*}
& 4x_1 - 4 + \lambda_1 - \lambda_2 = 0, \\
& 2x_2 - 4 + 2\lambda_1 + \lambda_2 = 0, \\
& \lambda_1 (x_1 + 2x_2 - 1) = 0, \\
& \lambda_2 (-x_1 + x_2) = 0, \\
& \lambda_1 \geq 0, \\
& \lambda_2 \geq 0.
\end{align*}
$$
Complementarity conditions yield 4 cases:

(i) If $\lambda_1 = 0$, $\lambda_2 = 0$, then $x_1 = 1$, $x_2 = 2$. This violates the second constraint.

(ii) If $\lambda_1 = 0$, $-x_1 + x_2 = 0$, then $x_1 = \frac{4}{3}$, $x_2 = \frac{4}{3}$, $\lambda_2 = \frac{4}{3}$. This violates the first constraint.

(iii) If $x_1 + 2x_2 - 1 = 0$, $\lambda_2 = 0$, then $x_1 = \frac{5}{9}$, $x_2 = \frac{2}{9}$, $\lambda_1 = \frac{16}{9}$. This is a feasible solution.

(iv) If $x_1 + 2x_2 - 1 = 0$, $-x_1 + x_2 = 0$, then $x_1 = \frac{1}{3}$, $x_2 = \frac{1}{3}$, $\lambda_1 = 2$, $\lambda_2 = -\frac{2}{3}$. This violates $\lambda_2 \geq 0$.

Since

$$\nabla_x^2 L(x, \lambda) = \begin{pmatrix} 4 & 0 \\ 0 & 2 \end{pmatrix}$$

is (constant and) positive definite, the point $x^* = \left( \frac{5}{9}, \frac{2}{9} \right)$, with $\lambda^* = \left( \frac{16}{9}, 0 \right)$, satisfies the first and second order KKT conditions.

Inspection show that it is a strict minimizer.
Assume now that the functions $f(x)$ and $c_i(x)$, $i = 1, \ldots, m$, are twice continuously differentiable, and let $x^* \in \mathbb{R}^n$, $\lambda^* \in \mathbb{R}^m$ be such that

\[ \nabla_x L(x^*, \lambda^*) = 0, \]
\[ c_i(x^*) = 0, \text{ if } i \in \mathcal{E}, \]
\[ c_i(x^*) \leq 0, \text{ if } i \in \mathcal{I}, \]
\[ \lambda_i^* c_i(x^*) = 0, \]
\[ \lambda_i^* \geq 0, \text{ for } i \in \mathcal{I}, \]

and

\[ d^\top \nabla_{xx}^2 L(x^*, \lambda^*) d > 0, \]

for all $d$, such that $\nabla c_i(x^*)^\top d = 0$, $i \in \mathcal{E} \cup \mathcal{A}(x^*)$.

Assume also that $\lambda_i^* > 0$ for all active inequality constraints (strict complementary slackness condition).

Then $x^*$ is a strict local minimizer of $f(x)$ subject to the constraints.
A practical way of solving constrained optimization problems is by adding a *penalty function* per constraint to the objective function.

These penalty functions attempt to mimic the effect of the constraint functions, i.e. discourage the optimizer to search outside of the feasible set.

Typically, one considers a sequence of such modified objective functions, in which the penalty functions are multiplied by larger and larger positive coefficients, thus forcing the optimizer to stay closer and closer to the feasible set.

Two commonly used penalty function methods are:

(i) *Exterior penalty functions* impose a penalty for violation of constraints.

(ii) *Interior penalty functions* impose a penalty for approaching the boundary of an inequality constraint.
The simplest exterior penalty method is the **quadratic penalty function**.

The quadratic penalty function applied to constraint optimization problems with equality constraints only, \( c_i(x) = 0 \), for \( i = 1, \ldots, m \), is defined as

\[
Q(x, \mu) = f(x) + \frac{\mu}{2} \sum_{i=1}^{m} c_i(x)^2,
\]

(32)

where \( \mu > 0 \) is the **penalty parameter**.

As \( \mu \to \infty \), the presence of the penalty term imposes more and more severe penalty for constraints violations.
Algorithmically, we consider a sequence $\mu_0 < \ldots < \mu_k \to \infty$, and exit criteria

$$\|\nabla Q(x, \mu)\|_\infty \leq \varepsilon_k,$$

(33)

where $\varepsilon_k > 0$, $k = 0, 1, \ldots$, is a sequence with $\varepsilon_k \to 0$, as $k \to \infty$.

We start the search with an initial guess $x_0$.

For $k = 0, 1, \ldots$, we find a minimizer $x^*_k$ of $Q(x, \mu_k)$ satisfying the stopping criterion (33).

In the next step, we increase $\mu_k$ to $\mu_{k+1}$, initialize the search with $x^*_k$ (“warm start”), and find a minimizer $x^*_{k+1}$ satisfying (33).

We stop, once the final convergence test is satisfied.

The expectation is that, as $k \to \infty$, $x^*_k$ converges to the minimizer of the original problem, and thus, for $\mu_k$ sufficiently large, $x^*_k$ will be a good approximation to $x^*$. 
Example 7. Consider again the problem:

$$\min x_1 + x_2, \quad \text{subject to } x_1^2 + x_2^2 - 2 = 0.$$ 

The quadratic penalty function is given by

$$Q(x, \mu) = x_1 + x_2 + \frac{\mu}{2} (x_1^2 + x_2^2 - 2)^2.$$ 

We find that the gradient and Hessian of $Q(x, \mu)$ are

$$\nabla Q(x, \mu) = \begin{pmatrix} 1 + \mu x_1(x_1^2 + x_2^2 - 2) \\ 1 + \mu x_2(x_1^2 + x_2^2 - 2) \end{pmatrix},$$

and

$$\nabla^2 Q(x, \mu) = \mu \begin{pmatrix} 3x_1^2 + x_2^2 - 2 & 2x_1x_2 \\ 2x_1x_2 & x_1^2 + 3x_2^2 - 2 \end{pmatrix},$$

respectively.
Example

We choose $\mu_0 = 1$ and perform Newton’s search with the initial guess $x_0 = (0, 0.1)$.

The table below shows the values of $x_k^*$ corresponding to $\mu_k = k$.

<table>
<thead>
<tr>
<th>$\mu_k$</th>
<th>$x_k^*$</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>(-1.19148788, -1.19148788)</td>
</tr>
<tr>
<td>2</td>
<td>(-1.10715987, -1.10715987)</td>
</tr>
<tr>
<td>3</td>
<td>(-1.07474445, -1.07474445)</td>
</tr>
<tr>
<td>4</td>
<td>(-1.05745377, -1.05745377)</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td>100</td>
<td>(-1.00249069, -1.00249069)</td>
</tr>
</tbody>
</table>
The sequence of $\mu_k$ can be chosen adaptively, based on the difficulty of minimizing $Q(x, \mu_k)$.

When minimizing $Q(x, \mu_k)$ turns out to be difficult, in the next step, we choose $\mu_{k+1}$ to be modestly larger than $\mu_k$.

Otherwise, we can take $\mu_{k+1}$ significantly larger than $\mu_k$.

As $\mu_k \to \infty$, the search proves in general to get increasingly harder due to the fact that the Hessian of $Q(x, \mu_k)$ becomes close to ill conditioned.

This phenomenon can be seen in Example 7, where the magnitude of the Hessian increases in proportion to $\mu$.

One can prove the following general theorem [2]: If $x_k^*$ is the exact global minimizer of $Q(x, \mu_k)$, then every limit point $x^*$ of $x_k^*$ is a global solution to the original problem.
The quadratic penalty function can be defined for a problem with both equality and inequality constraints as follows:

\[ Q(x, \mu) = f(x) + \frac{\mu}{2} \sum_{i \in E} c_i(x)^2 + \frac{\mu}{2} \sum_{i \in I} (c_i(x)^+)^2, \]  

(34)

where \( a^+ = \max(a, 0) \).
A refinement of the quadratic penalty method is the *augmented Lagrange method*. It reduces the possibility of ill conditioning discussed above by introducing approximate Lagrange multipliers into the objective function. We discuss the case equality constraints only, see [2] and [1] for the general case.

We consider the augmented Lagrange function

$$L_\mu(x, \lambda) = f(x) + \sum_{m=1}^{m} \lambda_i c_i(x) + \frac{\mu}{2} \sum_{i=1}^{m} c_i(x)^2.$$  \hspace{1cm} (35)

It differs from the standard Lagrange function (11) by the presence of the quadratic penalty term.

Note that the $\lambda_i$’s are not, strictly speaking, the Lagrange multipliers, as the critical points of $L_\mu(x, \lambda)$ are not feasible points for the original problem.
We do expect that as $\mu_k \to \infty$, the augmented Lagrange function will tend to the standard Lagrange function $L_{\mu_k}(x^*_k, \lambda) \to L(x^*, \lambda^*)$.

Heuristically, the first order condition reads:

$$0 \approx \nabla_x L_\mu(x, \lambda) = \nabla f(x) + \sum_{m=1}^{m} (\lambda_i + \mu c_i(x)) \nabla c_i(x),$$

from which we deduce that the “true” Lagrange multipliers are

$$\lambda_i^* \approx \lambda_i + \mu c_i(x),$$

or

$$c_i(x) \approx \frac{\lambda_i^* - \lambda_i}{\mu}.$$
In other words, if $\lambda_i$ is close to $\lambda_i^*$, the infeasibility of $x$ will be much smaller than $1/\mu$.

This suggests that the value of $\lambda_i$ at each iteration step $k$ (denoted by $\lambda_{i,k}$) is updated according to the following rule:

$$\lambda_{i,k+1} = \lambda_{i,k} + \mu_k c_i(x_k^*),$$  \hspace{1cm} (36)

for all $i = 1, \ldots, m$. 

A. Lesniewski

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Algorithmically, we consider a sequence $\mu_0 < \ldots < \mu_k \to \infty$, and stopping criteria

$$\|\nabla_x L_\mu(x, \lambda)\|_\infty \leq \varepsilon_k,$$

(37)

where $\varepsilon_k > 0$, $k = 0, 1, \ldots$, is a sequence with $\varepsilon_k \to 0$, as $k \to \infty$.

We start the search with an initial guess $x_0$ and $\lambda_0$.

For $k = 0, 1, \ldots$, we find a minimizer $x_k^*$ of $L_{\mu_k}(x, \lambda_k)$ satisfying the exit criterion (37).

In the next step, we increase $\mu_k$ to $\mu_{k+1}$, update $\lambda_k$ to $\lambda_{k+1}$ according to the rule (36), initialize the search with $x_k^*$, and find a minimizer $x_{k+1}^*$ satisfying (37).

We exit, once the final convergence test is satisfied.

The expectation is that, as $k \to \infty$, $x_k^*$ converges to the minimizer of the original problem without the necessity of sending $\mu_k \to \infty$. This should mitigate the danger of ill conditioning and speed up the convergence.
Consider again Example 7. The augmented Lagrange function reads

\[ L_\mu(x, \lambda) = x_1 + x_2 + \lambda(x_1^2 + x_2^2 - 2) + \frac{\mu}{2} (x_1^2 + x_2^2 - 2)^2. \]

We find that the gradient and Hessian of \( L_\mu(x, \lambda) \) are

\[ \nabla L_\mu(x, \lambda) = \begin{pmatrix} 1 + 2\lambda x_1 + \mu x_1 (x_1^2 + x_2^2 - 2) \\ 1 + 2\lambda x_2 + \mu x_2 (x_1^2 + x_2^2 - 2) \end{pmatrix}, \]

and

\[ \nabla^2 L_\mu(x, \lambda) = \begin{pmatrix} 2\lambda + \mu (3x_1^2 + x_2^2 - 2) & 2\mu x_1 x_2 \\ 2\mu x_1 x_2 & 2\lambda + \mu (x_1^2 + 3x_2^2 - 2) \end{pmatrix}, \]

respectively.
Example

We choose $\mu_0 = 1$ and perform Newton’s search with the initial guess $x_0 = (0, 0.1)$ and $\lambda_0 = 0.3$.

The table below shows the values of $x_k^*$ and $\lambda_k$ corresponding to $\mu_k = k$.

<table>
<thead>
<tr>
<th>$\mu_k$</th>
<th>$\lambda_k$</th>
<th>$x_k^*$</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>0.62706599</td>
<td>(-1.07867187, -1.07867187)</td>
</tr>
<tr>
<td>2</td>
<td>0.40227548</td>
<td>(-0.97149492, -0.97149492)</td>
</tr>
<tr>
<td>3</td>
<td>0.58299634</td>
<td>(-1.01494835, -1.01494835)</td>
</tr>
<tr>
<td>4</td>
<td>0.28520661</td>
<td>(-0.9901962, -0.9901962)</td>
</tr>
<tr>
<td>5</td>
<td>0.69491662</td>
<td>(-1.02027986, -1.02027986)</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td>10</td>
<td>0.31472584</td>
<td>(-0.99044963, -0.99044963)</td>
</tr>
</tbody>
</table>
Barrier methods

- Exterior penalty methods allow for generating infeasible points during the search. Therefore they not suitable when feasibility has to be strictly enforced.
- This could be the case if the objective function is undefined outside of the feasible set.
- Barrier methods are similar in spirit to the external penalty method.
- They generate a sequence of unconstrained modified differentiable objective functions whose unconstrained minimizers are expected to converge to the solution of the constrained problem in the limit.
Barrier methods belong to the category of interior penalty function methods and apply to inequality constraint optimization problems, i.e. \( \mathcal{E} = \emptyset \) in (1).

The essence of the method is to add a penalty term \( B(x) \) to the objective function, which has the following properties:

(i) It is defined and continuous whenever \( c_i(x) < 0 \), for all \( i = 1, \ldots, m \).
(ii) It goes to \( +\infty \), whenever \( c_j(x) \uparrow 0 \), for any \( i \).

Commonly used barrier functions are:

\[
B(x) = - \sum_{i=1}^{m} \log (-c_i(x)) \quad \text{(logarithmic)}, \tag{38}
\]

and

\[
B(x) = - \sum_{i=1}^{m} \frac{1}{c_i(x)} \quad \text{(inverse)}. \tag{39}
\]
Algorithmically, barrier methods work similarly to the penalty function methods. We select a sequence of positive numbers $\mu_0 > \ldots > \mu_k \to 0$, and consider a sequence of functions:

$$R(x, \mu_k) = f(x) + \mu_k B(x). \quad (40)$$

We find a local minimizer $x_k^*$, and use it as the initial guess for the next iteration. A general convergence theorem guarantees that any limit point of this sequence is a global minimizer of the original constrained problem.

Since the barrier function is defined only in the interior of the feasible set, any successive iteration must also be an interior point.

The barrier term $\mu_k B(x)$ goes to zero for interior points as $\mu_k \to 0$. The barrier term becomes thus increasingly irrelevant for interior points, while allowing $x_k^*$ to move closer to the boundary of $\Omega$.

This behavior is expected if the solution to the original constraint problem lies on the boundary.
Example 8. Consider the problem:

$$\min \frac{1}{2} (x_1^2 + x_2^2), \quad \text{subject to } x_1 \geq 2.$$ 

Its optimal solution is clearly $x^* = (2, 0)$.

The function

$$R(x, \mu_k) = \frac{1}{2} (x_1^2 + x_2^2) - \mu_k \log(x_1 - 2)$$

has a global minimum at

$$x_k^* = (1 + \sqrt{1 + \mu_k}, 0).$$

As $\mu_k$ decreases to zero, the unconstrained minimum $x_k^*$ converges to the constrained minimum $x^*$. 
References
